

INDIAN INSTITUTE OF TECHNOLOGY ROORKEE

NPTEL

NPTEL ONLINE CERTIFICATION COURSE

Nonlinear Programming

Lec-12

Geometric Programming-II

Dr. S. K. Gupta

Department of Mathematics

Indian Institute of Technology Roorkee

Hello friends, so welcome to lecture series on nonlinear programming. In the last lecture we have seen what geometric programming problems are, we have seen that there are some complicated nonlinear problems which maybe non-convex and how to solve such type of problems using arithmetic geometric mean inequality. Now in that lecture we have seen only unconstrained posynomials optimization.

I mean if we have a problem of posynomials without any constraint then how can we use arithmetic geometric inequality to solve those type of problems. Now in this lecture we will see that if we have a constraint problem, constraint posynomial optimization with equality constraint, we are focusing only on equality constraint in this lecture. Then how can we solve such type of problems so let us see.

So what is the formulation we have a geometric programming problem with equality constraints, so consider the case of minimizing an objective function which is the sum of posynomials subject to the equality constraints okay.

(Refer Slide Time: 01:30)

GP Problem with Equality Constraints

Consider the case of minimizing an objective function which is the sum of posynomials subject to the equality constraints. That is,

$$\begin{aligned} \text{Min } z &= f(x), \\ \text{subject to: } g_i(x) &= \sum_{r=1}^{P(i)} c_{ir} u_{ir}(x), \quad i = 1, 2, \dots, n \end{aligned}$$

where $P(i)$ denotes the number of terms in the i^{th} constraint and

$$u_{ir} = \prod_{j=1}^n (x_j)^{a_{ijr}}.$$

That is minimization of $f(x)$ subject to $g_i(x) = \sum_{r=1}^{P(i)} c_{ir} u_{ir}(x)$ for $i = 1, 2, \dots, n$, where $P(i)$ denotes the number of terms in the i^{th} constraint, c_{ir} and $u_{ir}(x)$ are running from 1 to n , where $P(i)$ denotes the number of terms in the i^{th} constraint, c_{ir} is the same as the posynomial term. Now let us discuss the same thing again, now we already know that sum of U_i where i is running from 1 to n is greater than or equal to product U_i where i is running from 1 to n , $U_i / \partial_i^{a_i}$ this we have discussed in the last class, this we obtain from the arithmetic geometry mean inequality.

(Refer Slide Time: 02:38)

$$\sum_{i=1}^n U_i \geq \prod_{i=1}^n \left(\frac{U_i}{\delta_i} \right)^{\delta_i}, \quad \delta_1 + \delta_2 + \dots + \delta_n = 1, \quad \delta_i > 0 \quad \forall i$$

Let $\delta_1 + \delta_2 + \dots + \delta_n = \lambda \Rightarrow \left(\frac{\delta_1}{\lambda} \right) + \left(\frac{\delta_2}{\lambda} \right) + \dots + \left(\frac{\delta_n}{\lambda} \right) = 1$

$$\sum_{i=1}^n U_i \geq \prod_{i=1}^n \left(\frac{U_i}{\delta_i/\lambda} \right)^{\delta_i/\lambda}$$

$$= \left(\frac{\lambda U_1}{\delta_1} \right)^{\frac{\delta_1}{\lambda}} \left(\frac{\lambda U_2}{\delta_2} \right)^{\frac{\delta_2}{\lambda}} \dots \left(\frac{\lambda U_n}{\delta_n} \right)^{\frac{\delta_n}{\lambda}}$$

$$\Rightarrow \left(\sum_{i=1}^n U_i \right)^\lambda \geq \lambda^{\delta_1 + \delta_2 + \dots + \delta_n} \left(\frac{U_1}{\delta_1} \right)^{\delta_1} \left(\frac{U_2}{\delta_2} \right)^{\delta_2} \dots \left(\frac{U_n}{\delta_n} \right)^{\delta_n}$$

$$= \lambda^\lambda \prod_{i=1}^n \left(\frac{U_i}{\delta_i} \right)^{\delta_i}$$

Where sum of δ is 1, and all δ_i are greater than 0. Now if sum of δ is not 1, suppose it is some λ . So how can we apply this inequality there okay. So suppose let us suppose sum of δ_i is some λ , if sum of δ_i is some λ so of course we cannot apply this inequality as much, because this is valid only when sum of δ is 1. So how can we apply this inequality, we divide each δ_i/λ to make the sum equal to 1.

So what we have got to do basically we write $\delta_1/\lambda + \delta_2/\lambda$ and so on $\delta_n/\lambda=1$, so that means to apply this inequality for this type of problems we replace $\delta_i/ \delta_i/\lambda$. Because now the new δ_i is δ/λ where sum is 1 okay, because sum of these, now this is one δ , this is second δ , this is third δ and another δ and sum of this is 1. And it is applicable only when the sum of δ is 1. So that means to apply this inequality for this problem we replace δ_i/λ , okay. So what we obtain it is summation I running from 1 to n $Q_i \geq \text{product I running from 1 to n, } (U_i / \delta_i/\lambda)^{\delta_i/\lambda}$. This inequality for which system this equality holds when sum of δ is λ , okay.

Now what is the right hand side let us see, you can simplify first it is $(\delta, \lambda_1, u_1)^{\delta_1/\lambda}$ the second term is $(\lambda, \delta_2, u_2)^{\delta_2/\lambda}$ and another term is $(\delta, \lambda_n, u_n)^{\delta_n/\lambda}$, okay. This is the right hands side now you first raised whole raise to the power λ both the side so this will implies summation I running from 1 to n, $(u_i)^\lambda$ in the left hand side when you do this so it is and you collect the powers of λ , okay. It is λ raise to the power from here when you take this λ to the left hand side and collect the powers of λ from each term.

So it is $\lambda^{\delta_1}, \lambda^{\delta_2}$ so it will add up another remaining term will be $(u_1/\lambda_1)^{\lambda_1} (u_2/\lambda_2)^{\lambda_2}$ and so $(u_n/\lambda_n)^{\lambda_n}$ so this sum is again λ from this expression we are solving this expression from this sum of δ is λ okay so this is equal to λ^λ and this is product $\prod_{i=1}^n (u_i/\lambda_i)^{\lambda_i}$ so we have obtained another equality that is if $u_1 + u_2$ and so on up to u_n whole raised to the power λ is always $\geq \lambda^\lambda$ product of $(u_i/\lambda_i)^{\lambda_i}$ where sum of δ is λ , okay.

So we will use this inequality for solving GP problem or geometric programming problem with equality type constraints. How we will do that let us see, let us discuss this by an example, okay.

(Refer Slide Time: 07:05)

Problem

- Min $f = 40x_1^{-1}x_2^{-1}x_3^{-1} + 40x_1x_3$
- subject to: $4x_1x_2 + 2x_2x_3 = 8,$
- $x_1, x_2, x_3 > 0.$

So to illustrate that how can we solve such type of problem using equality constraint let us discuss it by an example, suppose we have this example it is minimizing of f which is equal to $40x_1^{-1}x_2^{-1}x_3^{-1} + 40x_1x_3$ subject to you divide it by 8 this constraint you divide it by 8 so what we obtain it is $\frac{1}{2}x_1, x_2 + \frac{1}{4}, x_2x_3 = 1$ and x_1, x_2, x_3 are strictly greater than 0, okay.

(Refer Slide Time: 07:49)

$$\begin{aligned}
 \text{Min } f &= 40x_1^{-1}x_2^{-1}x_3^{-1} + 40x_1x_2 \\
 \text{s.t. } & \frac{1}{2}x_1x_2 + \frac{1}{4}x_2x_3 = 1, \\
 & x_1, x_2, x_3 > 0.
 \end{aligned}$$

$$\begin{aligned}
 \frac{1}{2}x_1x_2 + \frac{1}{4}x_2x_3 &= 1 \\
 u_3 + u_4 &= 1 \\
 1 &= 1^\lambda = (u_3 + u_4)^\lambda \\
 &\geq \lambda^\lambda \left(\frac{u_3}{\delta_3}\right)^{\delta_3} \left(\frac{u_4}{\delta_4}\right)^{\delta_4}, \quad \delta_3 + \delta_4 = \lambda, \\
 & \quad \delta_3, \delta_4 > 0.
 \end{aligned}$$

$$\begin{aligned}
 f &= 40x_1^{-1}x_2^{-1}x_3^{-1} + 40x_1x_2 \\
 &= u_1 + u_2 \\
 &\geq \left(\frac{u_1}{\delta_1}\right)^{\delta_1} \left(\frac{u_2}{\delta_2}\right)^{\delta_2}, \quad \delta_1 + \delta_2 = 1, \\
 & \quad \delta_1, \delta_2 > 0 \\
 &\geq \left(\frac{u_1}{\delta_1}\right)^{\delta_1} \left(\frac{u_2}{\delta_2}\right)^{\delta_2} \lambda^\lambda \left(\frac{u_3}{\delta_3}\right)^{\delta_3} \left(\frac{u_4}{\delta_4}\right)^{\delta_4}, \quad \delta_3 + \delta_4 = \lambda, \\
 & \quad \delta_3, \delta_4 > 0. \\
 &= \lambda^\lambda \left(\frac{40x_1^{-1}x_2^{-1}x_3^{-1}}{\delta_1}\right)^{\delta_1} \left(\frac{40x_1x_2}{\delta_2}\right)^{\delta_2} \left(\frac{\frac{1}{2}x_1x_2}{\delta_3}\right)^{\delta_3} \left(\frac{\frac{1}{4}x_2x_3}{\delta_4}\right)^{\delta_4}
 \end{aligned}$$

So first let us focus on the objective is $f = 40 x_1^{-1} x_2^{-1} + 40 x_1x_3$ now it is request to $u_1 + u_2$ okay, let u suppose a first term is u_1 the second term is u_2 so it is greater than equal to by the same inequality $(u_1/ \delta_1)^{\delta_1} (u_2/ \delta_2)^{\delta_2}$ were $\delta_1 + \delta_2$ were $\delta_1 + \delta_2$ is 1 and δ_1, δ_2 is strictly greater than 0. This is by the objective function, now come to the constraint what is the constraint we are having it is $\frac{1}{2} x_1 x_2 + \frac{1}{4} x_2 x_3 = 1$.

So it is u_3 suppose it is u_3 suppose it is $u_4 = 1$ were u_3 is the first over the constraint and u_4 is the second term over the constraint okay. Now to deal with this equation we discuss it like this one is always equals to 1 raise to the power λ for any λ okay so this one can be replaced by $u_3 + u_4$ is $1/\lambda$ because $u_3 + u_4$ is 1 so this one can be replaced by $u_3 + u_4$ and we just now discuss in the slide also that some of u_y is whole raise to power λ is always greater than $= \lambda^\lambda$.

(Refer Slide Time: 09:16)

Continued...

Therefore, the expression (1) gives

$$\begin{aligned}\sum_{i=1}^n u_i &\geq \prod_{i=1}^n \left(\frac{u_i}{\delta_i/\lambda} \right)^{\delta_i/\lambda} \\ \left(\sum_{i=1}^n u_i \right)^\lambda &\geq \prod_{i=1}^n \left(\frac{\lambda u_i}{\delta_i} \right)^{\delta_i} \\ &= \lambda^{(\delta_1 + \delta_2 + \dots + \delta_n)} \prod_{i=1}^n \left(\frac{u_i}{\delta_i} \right)^{\delta_i} = \lambda^\lambda \prod_{i=1}^n \left(\frac{u_i}{\delta_i} \right)^{\delta_i} \\ \Rightarrow \left(\sum_{i=1}^n u_i \right)^\lambda &\geq \lambda^\lambda \prod_{i=1}^n \left(\frac{u_i}{\delta_i} \right)^{\delta_i}\end{aligned}$$

In two product of $u_i / \delta_i^{\delta_i}$ where sum of δ_i is λ so this is $u_3 + u_4^\lambda$ so it will $> = \lambda^\lambda$ into $u_3 / \delta_3^{\delta_3}$
 $u_4 / \delta_4^{\delta_4}$ where $\delta_3 + \delta_4$ is λ and $\delta_3 \delta_4$ is r strictly written as 0, this is by in equality okay.

(Refer Slide Time: 10:00)

Problem

- Min $f = 40x_1^{-1}x_2^{-1}x_3^{-1} + 40x_1x_3$
subject to: $4x_1x_2 + 2x_2x_3 = 8$,
 $x_1, x_2, x_3 > 0$.



So what we obtained from here by the constraint we obtained this in equality now this we can use it here you see this we can always write this into 1 okay and this one is $1 \geq$ this in equality so this is $\geq u_1 / \delta_1^{\delta_1} u_2 / \delta_2^{\delta_2}$ into 1 and 1 is $\geq \lambda u_3 / \lambda_3^{\lambda_3}$ and $u_4 / \delta_4^{\delta_4}$ where $\lambda_3 + \lambda_4$ is λ λ_3 λ_4 are strictly > 0 okay so it is equals to λ^λ this can be put it outside the brackets okay what us u_1 , u_1 first sum of the objective function that is $40 s_1^{-1} x_2^{x_2-1} x_3^{-1}$ of δ_1 , what us u_2 , u_2 is $40 x_1 x_3 / \delta_2^{\delta_2}$ what is u_3 , u_3 is a third term.

First we constraint $x_1 x_2 / \delta^{\delta_3}$ and u_4 is $1/4 x_2 x_3 / \delta_4^{\delta_4}$ okay now again you will collect the powers of the variables here we are having 3 variables $x_1 x_2 x_3$ you collect the power of $x_1 x_2$ and x_3 put it equal to 0 and also $\delta_1 + \delta_2 = 1$ to the adjacent condition and try to find out the maximum value of the right hand side, so what are the equations we will obtained from this we can obtained.

(Refer Slide Time: 12:05)

$$f = 4x_1^2 x_2^2 x_3^2 + 4x_1 x_2$$

$$s.t. \quad \frac{1}{2} x_1 x_2 + \frac{1}{3} x_2 x_3 = 1$$

$$x_1, x_2, x_3 > 0$$

$$\frac{u_1}{x_1} = \frac{u_2}{x_2} \quad \frac{u_3}{x_3} = \frac{u_4}{x_4}$$

$$\Rightarrow \frac{40x_1^2 x_2^2 x_3^2}{u_1} = \frac{40x_1 x_2}{u_2} \quad \frac{\frac{1}{2} x_1 x_2}{\frac{1}{3}} = \frac{\frac{1}{3} x_2 x_3}{\frac{1}{3}} = \frac{2}{3}$$

$$\frac{u_1}{x_1} = \frac{u_2}{x_2} = \frac{u_3 + u_4}{x_1 + x_2} = \frac{1}{\frac{1}{2} + \frac{1}{3}} = \frac{2}{3}$$

$$\left. \begin{aligned} -\delta_1 + \delta_2 + \delta_3 &= 0 \\ -\delta_1 + \delta_3 + \delta_4 &= 0 \\ -\delta_1 + \delta_3 + \delta_4 &= 0 \\ \delta_1 + \delta_2 &= 1 \\ \delta_3 + \delta_4 &= \lambda \end{aligned} \right\} \Rightarrow \delta_1 = \frac{2}{3}$$

$$\delta_2 = \delta_3 = \delta_4 = \frac{1}{3}$$

$$\lambda = \delta_3 + \delta_4 = \frac{1}{3} + \frac{1}{3} = \frac{2}{3}$$

If you contain the powers of x_1 it is $-\delta_1 - \delta$ from 1 here $+\delta_2 + \delta_3 = 0$ $\delta_1 + \delta_2 + \delta_3 = 0$ for second term that is a power of x_2 it is $-\delta_1$ no x_2 is here so no power of δ_2 for x_2 and it is $\delta_3 + \delta_4$ that is $-\delta_1 + \delta_2 + \delta_4 = 0$ okay and the power of x_3 , x_3 is $-\delta_1 + \delta_2 - \delta_1 + \delta_2 / x_2 + \delta_2 + \delta_4$ power of x_1 is $-\delta_1 + \delta_2 + \delta_3 = 0 - \delta_1 + \delta_3 + \delta_4 = 0$ power of x_2 is -1 and $+\delta_3$ okay $+\delta_3 + \delta_4 = 0$ and $\delta_1 + \delta_2 = 1$ and $\delta_3 + \delta_4$ is λ .

So here we are having four equation four unknowns okay, we can solve it and find the values of δ_1, δ_2 and δ_3 and δ_4 of course okay. So how to now when we solve this equation what are values $\delta_1, \delta_2, \delta_3, \delta_4$ we can obtain so you can solve it I know the values of $\delta_1, \delta_2, \delta_3$ for this system δ_1 can be obtain as $2/3$ and $\delta_1 = \delta_2 = \delta_3$ are $1/3$ you can solve it and find these values of $\delta_1, \delta_2, \delta_3$ and δ_4 okay, you get simply verify that these are satisfying all the constraints all the equations okay.

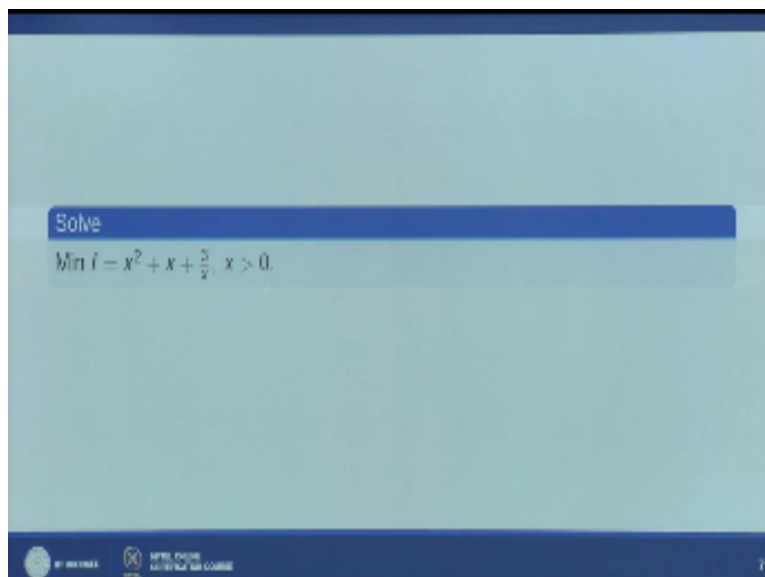
So what will be λ , λ is nothing but $\delta_3 + \delta_4$ which is okay, and $\delta_2 + \delta_3$ and δ_4 you can easily verify okay. Now $\delta_3 + \delta_4$ is $1/3 + 1/3$ that is $2/3$, so in this way we can find out the values of δ and λ . Now how we find out the value of the unknowns that x_1, x_2 and x_3 okay, so for that we will again apply the equality condition from the equality condition we obtain the equality hold everywhere for $\delta_1 = u_2 / \delta_2$ for the first one.

And $u_3 / \delta_3 = u_4 / \delta_4$ for the constraint also the equality will hold when this condition holds we can easily verify. Now from here we can obtain u_1 is $40 x_1^{-1} x_2^{-1} x_3^{-1} / \delta_1$ is $2/3$ u_2 is $40 x_1 x_3 / \delta_2$

this we are having three terms 1, 2, 3 how many unknowns 2 unknowns 3 – 2- 1 that is degree of difficulty 0.

Was this problem degree of difficulty is because we are having only three terms with one unknown so 3-1-1 that is one so degree of difficulty for this is one. Now it is easy it is I mean easy if a degree of difficulty is 0 because in that case we are having a unique solution if exist okay so the right hand side of the, an equality will be having a fixed value. Now if that degree of difficulty is more than or equal to one so how can we solve this type of problems. So for illustration let us discuss this problem.

(Refer Slide Time: 19:14)



We already know a degree of difficulty is this is one because it is having one unknown at three terms so 3-1-1 that is one, so for this problem it is again unconstraint problem this is simply illustration that $u_1 + u_2 + u_3 \geq u_1 / \delta_1^{\delta_1} u_2 / \delta_2^{\delta_2} u_3 / \delta_3^{\delta_3}$, we are sum out δ is one and δ I are

written 0 for all x . Now when you connect the power of x here only one variable is there, so how many equations we obtain only one.

And the second constraint is this and how many unknown's 3 unknown's $\delta_1, \delta_2, \delta_3$ okay so will be the equation it is where is substitute it here it is $x^2 / \delta_1^{\delta_1}$ it is $x / \delta_2^{\delta_2}$ it is $3 / \delta_3^{\delta_3}$, so when you connect the power of x put it equal to 0 we obtain $2\delta_1 + \delta_2 - \delta_3 = 0$. So here we are having only two equations with three unknowns so it will be having infinitely minor solutions, so how can we solve such problems.

Now what is our right hand side our right hand side it when we make it free from x it is $1 / \delta_1^{\delta_1} 1 / \delta_2^{\delta_2} 3 / \delta_3^{\delta_3}$ we want to find out the maximum value of this expression maximum value will give the mini value of F okay. Now what we will do we will try to express all the variables in terms of one variable say it δ_3 okay we will try to express that we can easily do you see $\delta_1 + \delta_2 + \delta_3$ is one and $\delta_2 \delta_1 + \delta_2 - \delta_3 = 0$ okay.

So you can subtract these two equations so it is $-\delta_1 + 2\delta_3 = 1$ so δ_1 will be nothing but $2\delta_3 - 1$ so that is in terms of δ_3 so you can substitute do this δ_3 over here or in the any one of the equation when we substituted over here suppose so δ_1 is $2\delta_3 - 1 + \delta_2 \delta_3 = 1$, so δ_2 we obtain as $2 - 3\delta_3$ okay it is δ_2 and it is one $2 - \delta_3$ okay, so we have express all the variables in terms of one variable that is δ_3 .

Now you can put it over here now you focus only on this term okay what is this term say it is $g \delta$, $g \delta$ is a right hand side $1 / \delta_1$, δ_1 is $2\delta_3 - 1$ okay. So in this way we have expressed in right hand side only in terms of one variable now what we have to do we have to find out the maximum value of this expression which is only one variable how we can do that you just take the logarithm for both the side you just take the log for both the sides find out the first derivative respective Δ_3 put it in equal to 0 that will give the maximum value of $g \Delta$ for which you can find out the value for Δ_3 for which $g \Delta$ is maximum once you find the value of Δ_3 for which $g \Delta$ is maximum you can substitute Δ_3 over here .

And you can find out the value of Δ_1 and Δ_2 and once you found out the value of Δ_1 and Δ_2 then you can easily say that this equality will hold when are all equal that is $u_1 / \Delta_1 = u_2 / \Delta_2 = u_3 / \Delta_3$ and that using that you can easily solve you can easily find out the value of x for which affix is minimum okay.

So in order to find out the maximum value of this you simply take log book both the sides differentiate respective Δ^3 put it in equal to 0 in the first derivative and you can find out the value of Δ_3 for which $g\Delta$ is maximum you can substitute that Δ_3 over here to find out the value of Δ_1 and Δ_2 .

And that will give the values of $\Delta_1\Delta_2$ and Δ_3 to find out the maximum value of minimum value of f you use the equality condition okay so in this way we can solve the problems with constraint normal optimal with equality constraint or if we have a problem having difficulty more than 0 so thank you very much.

For Further Details Contact
Coordinator, Educational Technology Cell
Indian Institute of Technology Roorkee
Roorkee – 247667
E Mail: etcell.iitrke@gmail.com, etcell@iitr.ernet.in
Website: www.iitr.ac.in/centers/ETC, www.nptel.ac.in

Camera
Jithin. K
Graphics
Binoy. V. P
Online & Video Editing
Mohan Raj. S

Production Team
Sarath Koovery
Arun. S
Pankaj Saini
Neetesh Kumar
Jitender Kumar
Nibedita Bisoyi

An Educational Technology cell
IIT Roorkee Production
© Copyright All Rights Reserved